

Module 6: CPU Scheduling

- Basic Concepts
- Scheduling Criteria
- Scheduling Algorithms
- Multiple-Processor Scheduling
- Real-Time Scheduling
- Algorithm Evaluation

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Basic Concepts

- Maximum CPU utilization obtained with multiprogramming
- CPU-I/O Burst Cycle – Process execution consists of a cycle of CPU execution and I/O wait.
- CPU burst distribution

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Alternating Sequence of CPU And I/O Bursts

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Histogram of CPU-burst Times

Burst Duration (μs)	Frequency
0	0
2	150
4	40
6	10
8	5
10	2
12	1
14	1
16	1
18	1
20	1
22	1
24	1
26	1
28	1
30	1
32	1
34	1
36	1
38	1
40	1

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CPU Scheduler

- Selects from among the processes in memory that are ready to execute, and allocates the CPU to one of them.
- CPU scheduling decisions may take place when a process:
 1. Switches from running to waiting state.
 2. Switches from running to ready state.
 3. Switches from waiting to ready.
 4. Terminates.
- Scheduling under 1 and 4 is *nonpreemptive*.
- All other scheduling is *preemptive*.

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Dispatcher

- Dispatcher module gives control of the CPU to the process selected by the short-term scheduler; this involves:
 - switching context
 - switching to user mode
 - jumping to the proper location in the user program to restart that program
- *Dispatch latency* – time it takes for the dispatcher to stop one process and start another running.

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Scheduling Criteria

- CPU utilization – keep the CPU as busy as possible
- Throughput – # of processes that complete their execution per time unit
- Turnaround time – amount of time to execute a particular process
- Waiting time – amount of time a process has been waiting in the ready queue
- Response time – amount of time it takes from when a request was submitted until the first response is produced, **not** output (for time-sharing environment)

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Optimization Criteria

- Max CPU utilization
- Max throughput
- Min turnaround time
- Min waiting time
- Min response time

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First-Come, First-Served (FCFS) Scheduling

- Example:

Process	Burst Time
P_1	24
P_2	3
P_3	3
- Suppose that the processes arrive in the order: P_1, P_2, P_3
The Gantt Chart for the schedule is:
- Waiting time for $P_1 = 0; P_2 = 24; P_3 = 27$
- Average waiting time: $(0 + 24 + 27)/3 = 17$

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FCFS Scheduling (Cont.)

Suppose that the processes arrive in the order P_2, P_3, P_1 .

- The Gantt chart for the schedule is:
- Waiting time for $P_1 = 6; P_2 = 0; P_3 = 3$
- Average waiting time: $(6 + 0 + 3)/3 = 3$
- Much better than previous case.
- *Convoy effect* short process behind long process

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Shortest-Job-First (SJF) Scheduling

- Associate with each process the length of its next CPU burst. Use these lengths to schedule the process with the shortest time.
- Two schemes:
 - nonpreemptive – once CPU given to the process it cannot be preempted until completes its CPU burst.
 - Preemptive – if a new process arrives with CPU burst length less than remaining time of current executing process, preempt. This scheme is known as the Shortest-Remaining-Time-First (SRTF).
- SJF is optimal – gives minimum average waiting time for a given set of processes.

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Example of Non-Preemptive SJF

Process	Arrival Time	Burst Time
P_1	0.0	7
P_2	2.0	4
P_3	4.0	1
P_4	5.0	4

- SJF (non-preemptive)
- Average waiting time = $(0 + 6 + 3 + 7)/4 = 4$

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Example of Preemptive SJF

Process	Arrival Time	Burst Time
P_1	0.0	7
P_2	2.0	4
P_3	4.0	1
P_4	5.0	4

- SJF (preemptive)

0 2 4 5 7 11 16

- Average waiting time = $(9 + 1 + 0 + 2)/4 = 3$

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Determining Length of Next CPU Burst

- Can only estimate the length.
- Can be done by using the length of previous CPU bursts, using exponential averaging.
 - t_n = actual length of n^{th} CPU burst
 - τ_{n+1} = predicted value for the next CPU burst
 - $\alpha, 0 \leq \alpha \leq 1$
 - Define:

$$\tau_{n+1} = \alpha t_n + (1 - \alpha)\tau_n$$

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Examples of Exponential Averaging

- $\alpha = 0$
 - $\tau_{n+1} = \tau_n$
 - Recent history does not count.
- $\alpha = 1$
 - $\tau_{n+1} = t_n$
 - Only the actual last CPU burst counts.
- If we expand the formula, we get:

$$\tau_{n+1} = \alpha t_n + (1 - \alpha) \alpha t_{n-1} + \dots + (1 - \alpha)^{n-1} \alpha t_1 + \dots + (1 - \alpha)^{n-1} t_0 \tau_0$$
- Since both α and $(1 - \alpha)$ are less than or equal to 1, each successive term has less weight than its predecessor.

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Priority Scheduling

- A priority number (integer) is associated with each process
- The CPU is allocated to the process with the highest priority (smallest integer = highest priority).
 - Preemptive
 - nonpreemptive
- SJF is a priority scheduling where priority is the predicted next CPU burst time.
- Problem \equiv Starvation – low priority processes may never execute.
- Solution \equiv Aging – as time progresses increase the priority of the process.

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Round Robin (RR)

- Each process gets a small unit of CPU time (*time quantum*), usually 10-100 milliseconds. After this time has elapsed, the process is preempted and added to the end of the ready queue.
- If there are n processes in the ready queue and the time quantum is q , then each process gets $1/n$ of the CPU time in chunks of at most q time units at once. No process waits more than $(n-1)q$ time units.
- Performance
 - q large \Rightarrow FIFO
 - q small \Rightarrow q must be large with respect to context switch, otherwise overhead is too high.

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Example: RR with Time Quantum = 20

Process	Burst Time
P_1	53
P_2	17
P_3	68
P_4	24

- The Gantt chart is:

0 20 37 57 77 97 117 121 134 154 162

- Typically, higher average turnaround than SJF, but better response.

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How a Smaller Time Quantum Increases Context Switches

process time = 10

quantum	context switches
10	0
5	1
1	9

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Turnaround Time Varies With The Time Quantum

process	time
P1	6
P2	3
P3	1
P4	7

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Multilevel Queue

- Ready queue is partitioned into separate queues: foreground (interactive) background (batch)
- Each queue has its own scheduling algorithm, foreground – RR background – FCFS
- Scheduling must be done between the queues.
 - Fixed priority scheduling; i.e., serve all from foreground then from background. Possibility of starvation.
 - Time slice – each queue gets a certain amount of CPU time which it can schedule amongst its processes; i.e., 80% to foreground in RR 20% to background in FCFS

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Multilevel Queue Scheduling

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Multilevel Feedback Queue

- A process can move between the various queues; aging can be implemented this way.
- Multilevel-feedback-queue scheduler defined by the following parameters:
 - number of queues
 - scheduling algorithms for each queue
 - method used to determine when to upgrade a process
 - method used to determine when to demote a process
 - method used to determine which queue a process will enter when that process needs service

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Multilevel Feedback Queues

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Example of Multilevel Feedback Queue

- Three queues:
 - Q_0 – time quantum 8 milliseconds
 - Q_1 – time quantum 16 milliseconds
 - Q_2 – FCFS
- Scheduling
 - A new job enters queue Q_0 which is served FCFS. When it gains CPU, job receives 8 milliseconds. If it does not finish in 8 milliseconds, job is moved to queue Q_1 .
 - At Q_1 job is again served FCFS and receives 16 additional milliseconds. If it still does not complete, it is preempted and moved to queue Q_2 .

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Multiple-Processor Scheduling

- CPU scheduling more complex when multiple CPUs are available.
- *Homogeneous processors* within a multiprocessor.
- *Load sharing*
- *Asymmetric multiprocessing* – only one processor accesses the system data structures, alleviating the need for data sharing.

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Real-Time Scheduling

- *Hard real-time* systems – required to complete a critical task within a guaranteed amount of time.
- *Soft real-time* computing – requires that critical processes receive priority over less fortunate ones.

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Dispatch Latency

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Algorithm Evaluation

- Deterministic modeling – takes a particular predetermined workload and defines the performance of each algorithm for that workload.
- Queuing models
- Implementation

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Evaluation of CPU Schedulers by Simulation

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